Note VIII

BY CHENCHAO DING

Dec. 2, 2024

```
type Name = String
data Expr where
    Var :: Name -> Expr
    Const :: Bool → Expr
    -- Rand :: Int -> Expr
    NotE :: Expr -> Expr
    OrE :: Expr -> Expr -> Expr
    AndE :: Expr -> Expr -> Expr
    PairE :: Expr -> Expr -> Expr
    Fst :: Expr -> Expr
    Snd :: Expr -> Expr
    LetE :: Name -> Expr -> Expr -> Expr
    deriving (Show)
data Value where
    Error :: Value
    BoolV :: Bool -> Value
    PairV :: Value -> Value -> Value
    deriving (Show)
type Outcome = Bool
type Env = [(Name, Value)]
type M = State Bool
```

```
lookupM :: Name -> Env -> M Value
lookupM s [] = return Error
lookupM s ((x,v):rest) = if (x == s) then (return v) else (lookupM s rest)
notM :: Value -> M Value
notM (BoolV b) = return (BoolV (not b))
notM = return Error
orM, andM :: Value -> Value -> M Value
orM (BoolV b1) (BoolV b2) = return (BoolV (or [b1, b2]))
orM = return Error
andM (BoolV b1) (BoolV b2) = return (BoolV (and [b1, b2]))
andM = return Error
interpM :: Env -> Expr -> M Value
interpM env (Var s) = lookupM s env
interpM env (Const b) = return (BoolV b)
interpM env (NotE e) = do
   v <- interpM env e
    notM v
```

```
pureContingent :: Expr -> M Outcome
pureContingent expr = get
contingentOnInterp :: Expr -> M Outcome
contingentOnInterp expr = do
   v <- interpM [] expr
   case v of
       BoolV b -> return b
      _ -> get
causeTrue :: Expr -> M Outcome
causeTrue expr = do
    put True
     contingentOnInterp expr
causeFalse :: Expr -> M Outcome
causeFalse expr = do
    put False
     contingentOnInterp expr
```

```
-- expression as quantum system

type Qsystem = Expr

-- test as single observable

type Observable = Qsystem -> M Outcome

-- test suite as measurement context

type Context = [Observable]

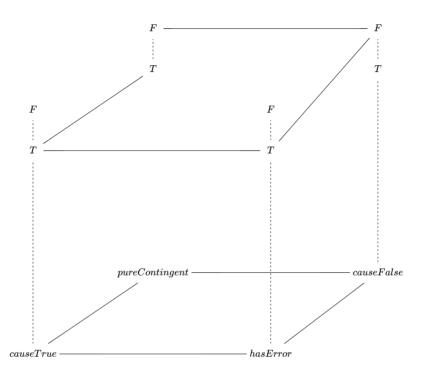
-- joint outcomes given a context

type Experiment = Qsystem -> Context -> M [Outcome]
```

Simulating Experiment

```
exp1 :: Qsystem -> Observable -> M Outcome
exp1 expr f = f expr
exp2 :: Qsystem -> (Observable, Observable) -> M (Outcome, Outcome)
exp2 expr (f1, f2) = do
   o1 <- f1 expr
   o2 <- f2 expr
   return (o1, o2)
     expn :: Qsystem -> Context -> M [Outcome]
     expn expr [] = return []
     expn expr (f:fs) = do
         o <- f expr
         os <- expn expr fs
         return (o:os)
type Seed = Int
-- joint outcomes given a context
runExperiment :: Qsystem -> Context -> Seed -> [Outcome]
runExperiment expr ctx seed =
    let m = expn expr ctx in
        evalState m (fst (randomBool seed))
```

```
-- Osystem independent contextuality...
     ctx1 :: [Expr -> M Outcome]
     ctx1 = [causeTrue, pureContingent]
     ctx2 :: [Expr -> M Outcome]
     ctx2 = [causeFalse, pureContingent]
     ctx3 :: [Expr -> M Outcome]
     ctx3 = [causeTrue, hasError]
     ctx4 :: [Expr -> M Outcome]
     ctx4 = [causeFalse, hasError]
     -- two by two (Bell configuration)
     suite1 :: [[Expr -> M Outcome]]
     suite1 = [ctx1,ctx2,ctx3,ctx4]
     -- one by one
     suite2 :: [[Expr -> M Outcome]]
     suite2 = [[pureContingent], [causeFalse], [hasError], [causeTrue]]
e1, e2, e3 :: Expr
e1 = NotE (Var "s")
e2 = LetE "x" (AndE (NotE (Const True)) (Const False)) (NotE (Var "x"))
e3 = PairE (PairE (Const True) e2) (Const True)
```



```
printResult :: Qsystem -> [Context] -> Seed -> IO ()
printResult expr [] seed = pure ()
printResult expr (c:cs) seed = do
    print (runExperiment expr c seed)
    printResult expr cs (seed + 1)
```

printResult e1 suite1 111

COMBINING CONTEXTUALITY AND CAUSALITY: A GAME SEMANTICS APPROACH

SAMSON ABRAMSKY, RUI SOARES BARBOSA, AND AMY SEARLE

ABSTRACT. We develop an approach to combining contextuality with causality, which is general enough to cover causal background structure, adaptive measurement-based quantum computation, and causal networks. The key idea is to view contextuality as arising from a game played between Experimenter and Nature, allowing for causal dependencies in the actions of both the Experimenter (choice of measurements) and Nature (choice of outcomes).

Definition 3 We say that a separated presheaf $A: C^{op} \to \mathbf{Sets}$ is no-signalling, or locally consistent, if every $A_{U \subseteq V}: A_V \to A_U :: s \mapsto s|_U$ is a surjection, i.e., if $A: C^{op} \to \mathbf{Surj}$ for the category \mathbf{Surj} of sets and surjections.